Networking in the Physical World

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Abstract

In this work we propose a network *meta-architecture* based on fundamental laws of physics and a physical model of computation. This meta-architecture may be used to frame discussions about novel network architectures as well as cross-layer alterations to the canonical network stack.

1 Introduction

The canonical layered network stack architecture has been leveraged successfully to create a world-wide computer network. However, technological advances have exposed limitations and weaknesses of the layered architecture. Large quantities of inexpensive computing power has encouraged software to encroach on the previously hardware-only physical layer. Quantum computing and quantum information theory have introduced new types of communication channels that do not necessarily fit cleanly into the existing stack architecture[1]. Much discussion has arisen around how to safely introduce cross-layer mechanisms into the stack[2], replace layers in the stack that have proven problematic in emerging applications, and even replacing the stack with something else entirely[3]. Unfortunately, most discussion surrounding cross-layer systems lacks a coherent and comprehensive set of goals: cross-layer dependencies are simply introduced as required to work around barriers in the existing network stack. Many "clean slate" networking architectures suffer from an excessively abstract point of view that is little help in actually constructing physically realizable and scalable systems. Conversely, approaches that make ad hoc modifications to the existing stack in the interest of functionality may not produce any real insight into the broader nature of networking systems. In this work we seek to rectify this situation by proposing a network *meta-architecture* that may be used to analyze and frame discussion about the canonical network stack and cross-layer modifications as well as clean slate architectures for the future. Our approach will begin with a model for computation, storage, and networking that is grounded

in the laws of physics as we currently know them. On top of this foundation we will then proceed to explore issues that are of importance to networking and storage systems.

2 A Physical Foundation for Computation

We will now outline a "bottom up" approach to abstract symbolic computation that directly relates abstract logical computation and physical operations. Instead of proposing a *symbol* manipulation machine we will begin with the manipulation and semantics of *physical states* and how to transform physical states into symbols. This model is based on fundamental tenets of modern physics, and was constructed under the influence of work by Feynman[4], Zuse[5] and Fredkin[6] regarding fundamental notions of computation and physics. We do not claim that this is a novel model, and this work is not primarily about computation. However, we feel that networking, storage, and computation are all fundamentally related tasks, and that having a firm notion of a computational model is critical for contextualizing and understanding networking and storage. Symbolic computation relies on physical entities. Abstract symbols map to physical entities and symbolic computation specifies changes in the physical states of those entities. Storage and networking move symbols by moving entities or properties of entities in time and space. Unless situated in some model of space and time, storage and networking are symbolically indistinguishable from the identity operator. Even systems that strive to eliminate physical reality introduce *some* model of location and time, *e.g.* memory addresses correspond to physical locations in a memory store, and the sequential execution of instructions marks the passing of time. IP addresses correspond to locations as well as names of systems, though this conflation of name and location causes well-documented problems when entities are allowed to move[7]. Counters used by logical clock systems express temporal event ordering and causality without insisting on non-local clock synchronization. Systems and applications operating directly in real space and time often use much more realistic models of space and time, *e.g.* the Global Positioning System (GPS) must account for relativistic and rotational motion effects on time and photons to obtain precise location estimates. We will refer to a physical entity as an *artifact*. An artifact has properties of physical particles or waves: location, mass, energy, momentum, charge, *etc.*, and quantum and relativistic effects may be important. An artifact may be a fundamental object, or assembled from other artifacts. Compound artifacts often behave differently from their subcomponents, introducing new levels of abstraction which are critical for managing complexity. For example, chemical analysis is facilitated by treating subatomic particles as aggregate atoms and molecules, and biology is greatly assisted by using cells as a basic abstraction.

An artifact may have attributes that correspond to some abstract unit of information, or *symbol*. When an artifact has such symbolic attributes we refer to that artifact as a *glyph*. Artifacts may interact with other artifacts. When an artifact alters another artifact we refer to the alterer as a *scribe*. Artifacts
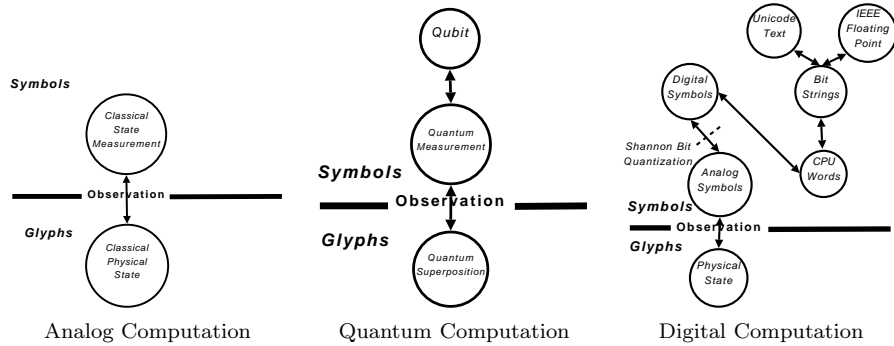
Figure 1: A diagram showing how glyphs, symbols, and filtering combine to create symbol sets for a analog, quantum, and traditional digital computation systems.

may be both glyphs and scribes; in fact a glyphs' action as a scribe is often used to observe it. *Physical operations* that alter artifacts may correspond to *symbolic operations*. A glyph must be *measured* to extract its corresponding symbol. This measurement is important and the subject of much engineering in classical computation systems. In quantum computation measurement, whether intentional or unintentional, is absolutely vital because the state superpositions upon which quantum computations rely collapse when observed. The measurement of a physical state includes some noise and uncertainty, ultimately limiting the number of symbols that may be reliably discerned from a particular physical state. Shannon's bit has emerged as the predominant way of handling noisy symbol streams, providing mathematical guarantees on the number of coded symbols that may be reliably embedded in a classical property of a glyph. Glyph systems used in digital computation systems may also be inherently quantized, providing an easy transition into bits so long as noise levels remain sufficiently low.

The basic mapping between glyphs and corresponding symbols is defined by the denotational semantics[8] of the information system, and additional layers of symbolic meaning may be defined by further denotations. We will not delve into a philosophical discussion of what the symbols *themselves* mean, or if it is meaningful to discuss the notion of a symbol outside of some physical embodiment. Instead we will simply state that the symbols may have further symbolic denotations defining their meanings and leave the rest to traditional denotational semantic techniques. Our only requirement is that any abstract symbol must eventually correspond to some physical embodiment. The operational and axiomatic constraints of glyphs are defined by the physical laws governing the glyphs, and these glyph semantics will "bubble up" and influence feasible symbolic operational[9] and axiomatic[10] semantics.

Figure 1 shows diagrams of how symbols for analog, quantum, and digital computation system are formed from glyphs. In the digital system, glyphs are measured to create analog symbols, which are then further processed to generate digital symbols, and then mapped into bit strings of machine word size.
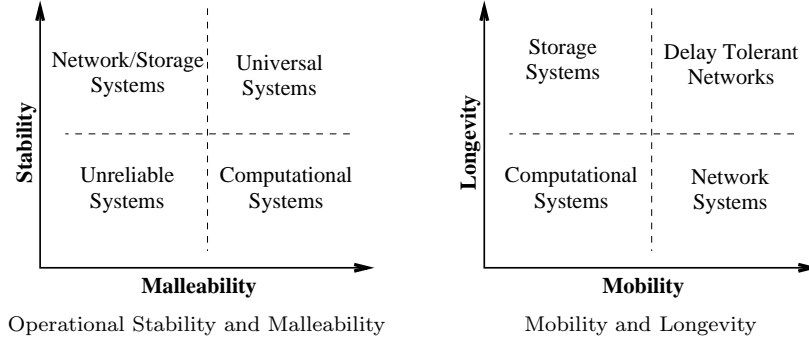
3

Figure 2: Computational, networking and storage systems place different demands on glyphs. Stability and malleability measure the difficulty and speed of glyph alteration respectively. Mobility and longevity measure how difficult it is to move the glyph in space and time respectively.

These machine words may be further mapped into operators of a microprocessor instruction set, Unicode characters of human language, or any number of other symbolic domains. Unfortunately the term "bit" is hopelessly overloaded: a "bit" not only refers to a specific way of converting a noisy signal into discrete symbols, but also the use of binary digits as a universal symbolic alphabet. When working with physically discretized systems it is convenient to use symbolic bits to represent states, but this is separate from performing bit-style signal filtering and discretizing. Current digital systems use multiple glyph repositories for different tasks. CPU registers are used for rapid, deterministic symbol manipulation embodied in a small number of glyphs, cache memory near the CPU registers stores a larger number of glyphs for rapid copying into registers, and external RAM and disk-based swap farther away holding an even greater number of glyphs. Typically all of these glyph stores are managed as a single virtual space that is accessed via sequential labels. Different glyph and symbol systems may have different sizes, potentially requiring padding or *fragmentation/reassembly* to match a symbolic system to a particular type of glyph. Quantum computing systems operate primarily on pure glyphs, making a symbolic observation only at points when collapsing superimposed quantum states has a very high chance of producing the desired computational result.

3   Physical System Requirements

Choosing an embodiment for networking, storage, or computation is fundamentally motivated by physical system characteristics. In traditional computation systems it is typical to use fast-moving bosons or a media-based waves for networking, and easily-trapped fermions for computation and storage. However, this choice is not universally true. As a microscopic example, consider the general operation of a cell. The normal operational "computation" is carried out

4

largely by proteins. These proteins are synthesized by cellular machinery driven by RNA, which in turn is synthesized from DNA stored in the nucleus. Tightly wound DNA is very compact and stable, making it an excellent medium for storage, but cannot be used directly by cell machinery to synthesize proteins. Transcribing segments of DNA into a more reactive and expanded form (RNA) allows the cell to synthesize the amino acids and proteins that drive cellular operation. As further examples, chemical and molecular communication schemes have been proposed for nanomachines[11], and the *Voyager* space probes contain discs intended to communicate information to any extraterrestrial entities that might stumble across them[12], likely at a very great distance in space and time. From a general system architecture design standpoint we propose four critical physical properties: *stability*, *i.e.* the difficulty/energy required to change a state, *malleability*, *i.e.* how quickly it may be changed, *longevity*, *i.e.* how easily it may be moved through time, and *mobility*, *i.e.* how easily it may be moved through space. These properties are illustrated in Figure 2. Expressing these properties quantitatively relies on an ability to handle physical space, time, and energy. Because networking and storage systems strive to preserve symbolic state we are also interested in *distortion*, *i.e.* unintended changes in glyph and/or symbolic state. In the next two sections we will discuss space, time, energy, and distortion in greater detail.

3.1   Space, Time, and Energy

Space, time, and energy are three obvious pieces of physical information that may be abstracted out of a physical system. In fact, these physical properties emerge throughout networking systems as well as physically-oriented applications, *e.g.* sensing and actuating systems. Current systems have varying notions of time, particularly when accuracy and precision greater than that specified by the POSIX *gettimeofday* call. Spatial location is largely *ad hoc*, but a latitude and longitude are widely-used as a coordinate system, largely driven by the increasing availability of GPS receivers. Naturally, latitude and longitude are not the end of the story: the DARPA XNAV[13] program proposes the use of x-ray pulsars for determining location and time in deep space, and a nano-scale machine would fare better with a correspondingly small coordinate system. Furthermore, space and time are inextricably linked. Most existing computational systems ignore relativistic factors, but as clocks become more and more precise and operating environments more and more extreme system designers will be forced to cope with modern physics, *e.g.* the GPS system itself requires accounting for relativistic effects in order to provide accurate location and time information.

3.2   Distortion

Distortion is perhaps not as obviously a fundamental physical concern as space, time, and energy, but distortion is critical for information systems. In this work we define distortion as an observable deviation from some "normal" state. Dis-

tortion is critical in networking and storage systems because networking and storage strive to produce changes in place and time *without* altering symbolic meaning, and distortions in glyphs may result in distortions in symbols. The effects of distortions have been extensively studied as noise in signal theory, Shannon's bit is a technique for handling distortions (*i.e.* noise) in an analyzable and generic fashion in quantized systems. Distortion generally only shows up at the lower layers in the canonical network stack because of Shannon's results regarding the separation of source and channel coding, but our meta-architecture does not require this split, *e.g.* for the exceptions to source/channel separation[14]. A more simplistic reason to permit violation of source and channel separation is the *cost* of performing encoding and decoding. The time and energy required for encoding and decoding in an information-theoretically optimal fashion may result in an overall suboptimal or inadequate computational system, and generic networking and storage systems should permit the computational portion of the system to make such cost/benefit decisions.

4   Names, Addresses, Locations, Times, and Routes

Earlier work[15] has stressed the fundamental nature of names, addresses, and routes. We will take a physical approach to names and addresses, starting with the notion that location and time are inherent properties of artifacts that may be observed and interpreted. *Names* are symbols associated with artifacts based on properties of those artifacts. We will refer to names that depend directly or indirectly on location and/or time as *addresses*. An address that depends exclusively on location and/or time may be used as part of a generic coordinate system for describing locations and times: we will refer to such names as *pure addresses*. We call location-independent names *labels*. Because we are treating location and time as universal artifact properties a label, by defining artifact(s), may define a set of addresses, and conversely addresses may define some set of labels. This fundamental connection between addresses and names can lead to confusion. For example, we consider an IP "address" to be a label because it is not defined by physical location. IP labels may be correlated with terrestrial location, but this is a side effect of how the labels have been allocated, *not* a fundamental property of the labels themselves. At a low level, addressing (coordinate) systems require reference points. An addressing scheme may have reference points defined in terms of other addressing schemes, and for abstract systems this process may continue indefinitely. However, for physical systems the process must be bootstrapped by an addressing scheme that does not depend on some other addressing scheme, *i.e.* in terms of some fundamental rigid body that is identifiable without referring to location. A *label* may be used to identify an origin that is external to a computational device, but the location of that origin must still be determined with respect to each computational device. We are left with the idea that each computational device must start with the ability to refer to "self" as a label and "here" as a location. The next logical step is for a system to have a model of itself

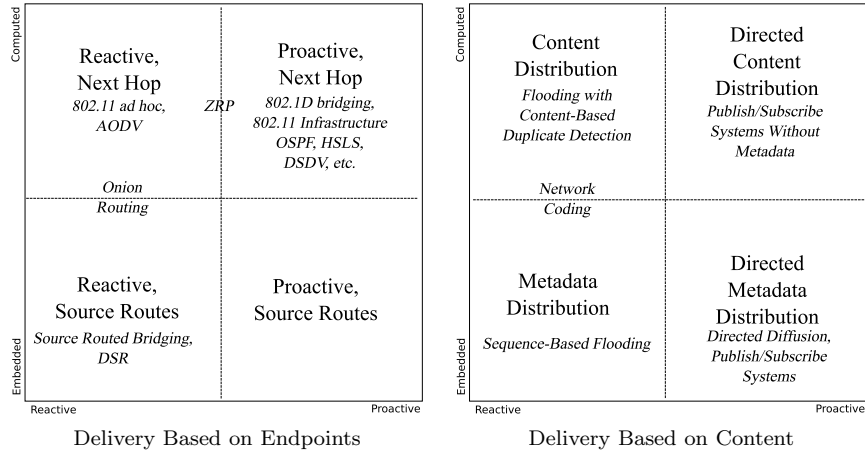|	Computed		Computed	
	Reactive, Next Hop *802.11 ad hoc, AODV*	*ZRP* — **Proactive, Next Hop** *802.1D bridging, 802.11 Infrastructure OSPF, HSLS, DSDV, etc.*	**Content Distribution** *Flooding with Content-Based Duplicate Detection*	**Directed Content Distribution** *Publish/Subscribe Systems Without Metadata*
	Onion Routing		*Network Coding*	
Embedded	**Reactive, Source Routes** *Source Routed Bridging, DSR*	**Proactive, Source Routes**	**Metadata Distribution** *Sequence-Based Flooding*	**Directed Metadata Distribution** *Directed Diffusion, Publish/Subscribe Systems*
	Reactive — Proactive		Reactive — Proactive	
	Delivery Based on Endpoints		**Delivery Based on Content**	

Figure 3: Route information may be embedded alongside payload or computed along the way. Routing decisions may be based on the content itself or on the endpoints to visit, and routes may be preemptively discovered and maintained or discovered and maintained on demand.

situated in space and time that may then be used to interact with the world. To this end, we propose that physical time and location as well as energy be treated as core system abstractions, and that tracking time and location as core system services.. Physical time and location may then be used to build shared addressing systems and to map between labels and addresses; the addition of energy permits stability, malleability, longevity, and mobility to be expressed. Multiple computational devices may define common addressing systems based on shared labels and artifacts that may then be used to construct shared maps of computational artifacts. Furthermore, physical time and location may be used for sensing

Core problems for networking systems are maintaining shared world maps and determining the routes and paths that are to be taken by glyphs. We will now broadly categorize different techniques that may be employed for routing. Routes may be precomputed and *embedded* with glyphs, as with source routes, or *computed* locally as glyphs travel, as with next hop routing. Routes may be calculated *proactively* ahead of time or *reactively* on request. Routes may be determined by *endpoints* or the symbols and glyphs (*content*). Endpoint routing is commonly used in existing topology-based routing schemes. Content-based routing does not track endpoint topology. Instead, glyph semantics [16] or cleverly merging symbols or glyphs[17] may be used to manage the flow of glyphs. Figure 3 illustrates how a selection of existing network routing protocols manage the motion of information between endpoints. Our categorization captures somewhat different aspects of routing than the categorization proposed by Hauzeur[15]. Our categorization does not capture the localization or frequency of route update calculations, but rather is concerned with what aspects of a system require labeling and addressing to function and how data

and metadata must be managed in a system. In most traditional layered digital systems metadata is kept at the "front edge" of the glyph and completely separate from the data. This arrangement of data and metadata has proven convenient, but is by no means the only possible arrangement. For example, the *Perspecta*[18] movie audio system contained sub-audible routing metadata for an audio stream alongside the audible soundtrack itself, and optical networks may opt to use light frequency to encode routing information. A "heap" structure for packet metadata has also been proposed [3], though we are not aware of any deployed systems built around this principle. Chemical and molecular networking schemes proposed for nanomachines[11] may use elements of both content and endpoint routing, and in general we feel that some combination of techniques should be supported to handle a broad range of environments. A further consideration for routing is *flow control*. Realistically there are limits on the number of glyphs that may occupy a single region of space at a given time, and flow control ensures that glyphs do not overfill containers. In the canonical network stack flow control is performed by MAC layers at individual links and over multiple links by higher layer protocols such as TCP. Applications may also perform their own flow control based on how quickly they can process and generate information.

5 System Meta-Architecture

In this section we will outline a networking system meta-architecture. In formulating this design we have drawn on a wealth of prior architectural systems work at the application layer[19], internetworking layer[20, 21], MAC and physical layers[22, 23, 24], and work spanning layers[25, 26, 27] for inspiration. Our major meta-architectural goals are fairly straightforward to state: we desire the ability to design system architectures that are *scalable*, *optimal*, *efficient*, *modular*, and *generic*. Furthermore, these system architectures must be feasible, though feasibility should be judged by the laws of physics and not necessarily current engineering limitations, and perform the tasks required by a networking system. We have already identified three primary tasks that networking must accomplish: labeling and addressing, determining paths in space and time for packets (routing and flow control), and fragmentation/reassembly to match desired symbol size to actual glyph capacity. We have also identified four basic physical properties that are common to networking systems: space, time, energy, and distortion. Three of these properties (space, time, and energy) are also of concern to physically-situated applications, even in the absence of a networking or storage system, and are candidates for general abstraction by the operating system for *all* sensor/actuator applications. Finally, we have described how generic computation, networking, and storage systems may be described in terms of semantic layers that ultimately reside within some physical artifact. With these goals, tasks, and fundamental physical information properties in hand we can now proceed to describe a layered meta-architecture in greater detail. Figure 4 illustrates our proposed meta-architecture as well as a decom-
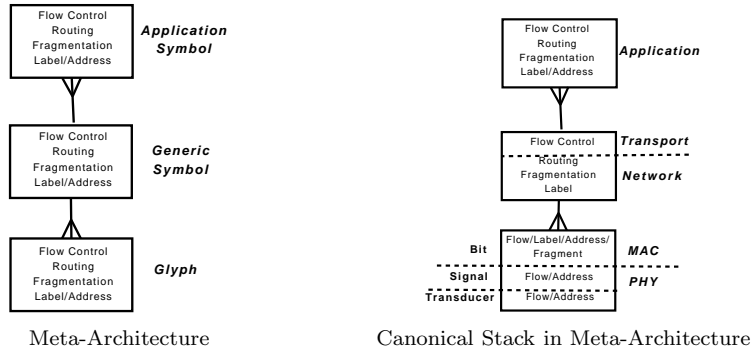
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| Meta-Architecture | Canonical Stack in Meta-Architecture |

Figure 4: A picture of our proposed network meta-architecture. The "crow's feet" connectors imply a potential one-to-many relationship, and the inside of each box enumerates the networking functions performed in each layer. The canonical stack shown also illustrates how Shannon's bit may be used to provide another set of layers based on measured signals and signal processing, forming another "waistline" at the bit transition. The canonical Software Defined Radio architecture[27] is built around the Shannon waistline and permits the entire gamut of network functions to be performed at the glyph-aware layer. Corresponding waistlines may also exist on the top end of the stack, such as the widespread use of XML as an interchange format for expressing application data semantics.

position of the canonical network stack. We will now briefly describe each layer of our meta-architecture. We will not delve into many interesting potential architectural details in this work, though we are actively exploring such details and plan on reporting them in future work.

5.1   Meta-Architecture Layers

The Glyph Layer   The *Glyph* layer is aware and has some model of the actual glyphs that will be used to transport symbols. An individual glyph layer may be aware of only a single type of glyph, though one could conceive of systems that provided "fast path" transitions between different types of glyphs or combine multiple glyph types into a single compound glyph. The Glyph layer is essentially a sensor/actuator system and will likely encounter problems similar to those found in sensor systems. This fact makes it seem that access to the physical world should be controlled and consolidated by the operating system at a level *below* the networking system to avoid a "stack inversion" when a sensor application attempts to use a networking system. Space, time, energy, and distortion at the Glyph layer are likely to be expressed in terms that are directly relevant to the glyph in question. For example, a directional wireless system has a "native" coordinate system based on its antenna pattern and electromagnetic wave propagation. Distortion may be expressed directly in terms of relevant glyph attributes. Using Shannon bits and signal processing provides an enormous boost in abstraction: signal-to-noise ratios allow systems to model links and perform power control in the context of control theory.

Routing and flow control at the Glyph layer have been extensively studied

in the context of MAC protocols, though most MAC layers do not perform multihop routing. Flow control may coordinate with the Generic Symbol layer by "backpressure" and routes may use whatever time, energy, space, and distortion metrics are sensible in a particular architecture. Connecting to the generic symbol layer also involves converting the transport glyph into a corresponding computational glyph, and details of this conversion will vary depending on the system in question. Using Shannon bits permits the formation of an architectural waistline at the Glyph layer that is not apparent from our generic meta-architectural diagram.

The Generic Symbol Layer   The Generic Symbol Layer is where we expect primary architectural waistlines to form, as is evidenced by the fanouts that exist on either side of it. This layer is essentially an "internetworking" layer that funnels in application requests for communication and fans them out to Glyph layer devices for actual transport. The Generic Symbol layer takes energy, time, and required address/label delivery constraints and determines the best use of intermediate Glyph devices based on their energy, time, distortion, and address/label delivery capabilities. This layer will likely face some of the greatest computational complexity as it must contend with potential combinatorial explosions on either end. This layer also seems a likely candidate for division into sublayers or some other more complicated systemic construct, though it should be stressed that having a non-layered *internal* structure of a layer in no way violates overall system layering.

The Application Symbol Layer   The Application Symbol layer is where individual applications may make networking and storage requests based on application-specific requirements. These requirements are generally driven by clients of the applications, and requests may be specified at a reasonably high level, *e.g.* discovering labels and addresses of other computational entities and exchanging information with them. Further layering of symbolic semantics may occur within the Application Symbol layer, *e.g.* the little-used OSI *Presentation* layer and the increasing use of XML as a standard data interchange format.

6   Conclusion

In this work we have outlined a physically grounded way of thinking about networking, computation, and storage and formulated a networking system *meta-architecture* in this context. We believe that such a meta-architecture is essential for organized analysis and discussion of network architectures and systems, and are actively exploring details of next-generation network architectures within the context of this meta-architecture.

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